Query Processing

Lecture topics:

- query interpretation
- basic operations
- costs of basic operations

References:

- text 3rd edition: chapter 18, sections 2-4, 7
- text 4th edition: chapter 15, sections 1–4, 8

Steps in query interpretation

- Translation
  - check SQL syntax
  - check existence of relations and attributes
  - replace views by their definitions
  - transform query into an internal form (similar to relational algebra)

- Optimization
  - generate alternative access plans, i.e., procedure, for processing the query
  - select an efficient access plan

- Processing
  - execute the access plan

- Data Delivery

Example

```
select V.Vno, Vname, count(*), sum(Amount)
from Vendor V, Transaction T
where V.Vno = T.Vno and V.Vno between 1000
   and 2000
group by V.Vno, Vname
having sum(Amount) > 100
```

- The following is a possible access plan for this query:
  - Scan the Vendor table, select all tuples where Vno is between 1000 and 2000, eliminate attributes other than Vno and Vname, and place the result in a temporary relation R_1
  - Join the tables R_1 and Transaction, eliminate attributes other than Vno, Vname, and Amount, and place the result in a temporary relation R_2. This may involve:
    - sorting R_1 on Vno
    - sorting Transaction on Vno
    - merging the two sorted relations to produce R_2

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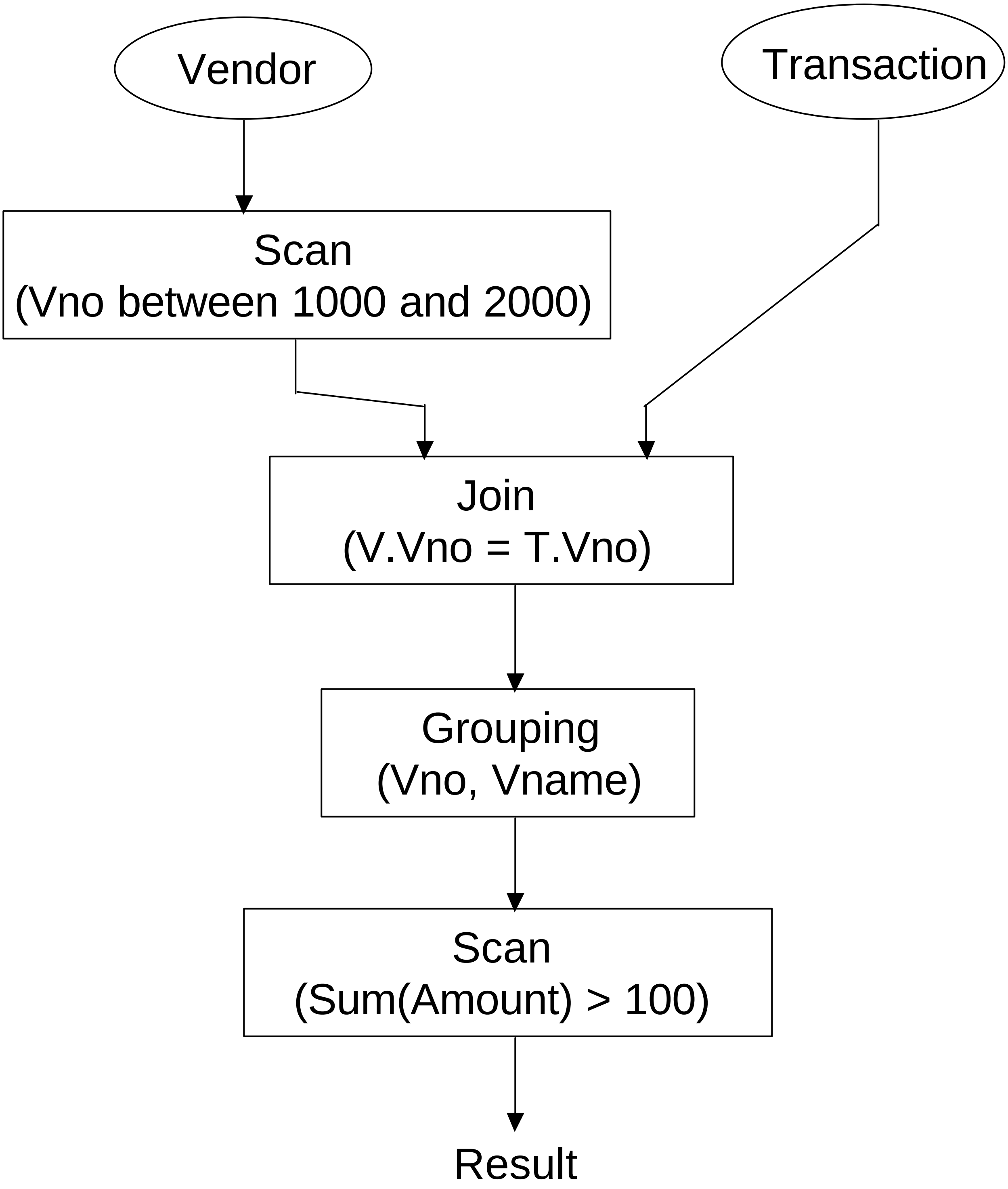
- Perform grouping on R_2, and place the result in a temporary relation R_3. This may involve:
  - sorting R_2 on Vno and Vname
  - grouping tuples with identical values of Vno and Vname
  - counting the number of tuples in each group, and adding their Amounts

- Scan R_3, select all tuples with **sum**(Amount) > 100 to produce the result.

Pictorial Access Plan



Some basic query processing operations

- Tuple Selection
  - without an index
  - with a clustered index
  - with an unclustered index
  - with multiple indices
- Projection
- Joining
  - nested loop join
  - hash join
  - sort-merge join
  - and others...
- Grouping and Duplicate Elimination
  - by sorting
  - by hashing
- Sorting

Tuple Selection

- sequential scanning:
  - when no appropriate index is available, we must check all file blocks holding tuples for the relation.
  - even if an index is available, scanning the entire relation may be faster in certain circumstances:
    - the relation is very small
    - the relation is large, but we expect most of the tuples in the relation to satisfy the selection criteria
  - E.g.:

```
select Anum
from Accounts
where Balance > 1.00
```

Joining relations

```
select C.Cnum, Balance
from Customer C, Accounts A
where C.Cnum = A.Cnum
```

- Two relations can be joined using the “nested loop” join method. Conceptually, that method works like this:

```
for each tuple c in Customer do
    for each tuple a in Accounts do
        if c.Cnum = a.Cnum then
            output c.Cnum,a.Balance
    end
end
```

Other techniques for join

- There are many other ways to perform a join. For example, if there is an index on the Cnum attribute of the Accounts relation, an **index join** could be used. Conceptually, that would work as follows:

```
for each tuple c in Customer do
        use the index to find Accounts tuples a
                where a.Cnum matches c.Cnum
        if there are any such tuples a then
                output c.Cnum, a.Balance
        end
end
```

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- other join techniques:
  - **sort-merge join**: sort the tuples of Accounts and of Customer on their Cnum values, then merge the sorted relations.
  - **hash join**: assign each tuple of Accounts and of Customer to a “bucket” by applying a hash function to its Cnum value. Within each bucket, look for Account and Customer tuples that have matching Cnum values.

Costs of basic operations

- Alternative access plans may be compared according to cost.
- The cost of an access plan is the sum of the costs of its component operations.
- There are many possible cost metrics. However, most metrics reflect the amounts of system resources consumed by the access plan. System resources may include:
  - disk block I/O’s
  - processing time
  - network bandwidth

A simple cost model

- The cost of an operation will be an estimate of the number of file block retrievals required to perform that operation.
- The cost of an operation will depend on the number of input tuples for that operation, and on the tuple blocking factor.
- The blocking factor is the number of tuples that are stored in each file block.
- For a relation *R*:
  - |*R*| will represent the number of tuples in *R*
  - *b*(*R*) will represent the blocking factor for *R*
  - |*R*|/*b*(*R*) is the number of blocks used to store *R*
- We will assume that an index search has a cost of 2.

Cost of selection

- relation Mark = (Studnum, Course, Assignnum, Mark)

```
select Studnum, Mark
from Mark
where Course = PHYS
 and Studnum = 100
 and Mark > 90
```

- Indices:
  - clustering index CourseInd on Course
  - non-clustering index StudnumInd on Studnum
- Assume:
  - |Mark| = 10000
  - *b*(Mark) = 50
  - 500 different students
  - 100 different courses
  - 10 different assignments

Selection Strategy: use CourseInd

- Assuming *uniform distribution* of tuples over the courses, there will be about |Mark|/100 = 100 tuples with Course = PHYS.

- Searching the CourseInd index has a cost of 2. Retrieval of the 100 matching tuples adds a cost of 100/*b*(Mark) data blocks, for a total cost of 4.

- Selection of *N* tuples from relation *R* using a clustered index has a cost of 2 + *N*/*b*(*R*)

Selection strategy: use StudnumInd

- Assuming *uniform distribution* of tuples over student numbers, there will be about |Mark|/500 = 20 tuples for each student.
- Searching the StudnumInd has a cost of 2. Since this is not a clustered index, we will make the pessimistic assumption that each matching record is on a separate data block, i.e., 20 blocks will need to be read. The total cost is 22.
- Selection of *N* tuples from relation *R* using a non-clustered index has a cost of 2 + *N*

Selection strategy: scan the relation

- The relation occupies 10,000/50 = 200 blocks, so 200 block I/O operations will be required.

- Selection of *N* tuples from relation *R* by scanning the entire relation has a cost of |*R*|/*b*(*R*)

Cost of joining

- relation Mark = (Studnum, Course, Assignnum, Mark)

- relation Student = (Studnum, Surname, Initials, Address, Birthdate, Sex)

-
```
select   Surname, Course,
         Assignnum, Mark
from Mark natural join Student
```

- Assume:
  - |Mark| = 10,000
  - *b*(Mark) = 50
  - |Student| = 1,000
  - *b*(Student) = 20

Cost of nested-loop join

```
for each Mark block M do
  for each Student block S do
    for each tuple m in M do
      for each tuple s in S do
        if m[Studnum] = s[Studnum] then
          join {m} and {s} and add to the result
```

- Each of the 200 blocks of Mark must be retrieved once.
- Each of the 50 blocks of Student must be retrieved once for each block of Mark, for a total of 200 x 50 = 10,000 student blocks retrieved.
- The total cost is 10,200 block retrievals.

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- The cost of a nested-loop join of R_1 and R_2 is either:

$$\frac{|R_1|}{b(R_1)} + \frac{|R_1|}{b(R_1)}\frac{|R_2|}{b(R_2)}$$

or

$$\frac{|R_2|}{b(R_2)} + \frac{|R_1|}{b(R_1)}\frac{|R_2|}{b(R_2)}$$

depending on which relation is the **outer** relation.

Cost of index join

- Suppose there is a unique, non-clustered index on Student [Studnum]

```
for each Mark block Bm do
  for each tuple m in Bm do
    use the index to locate Student tuples
    s with s[Studnum] = m[Studnum]
    join {m} and {s} and add to the result
```

- Each of Mark’s 200 blocks must be retrieved once. For each tuple of Mark, a single tuple of Student will be retrieved. This will require an index lookup and a data block retrieval, for a cost of 3.
- The total cost is 10000 * 3 + 200 = 30,200 block retrievals.

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- The cost of a unique index join, where R_1 is the outer relation, and R_2 is the inner, indexed relation, is:

$$\frac{|R_1|}{b(R_1)} + 3 \cdot |R_1|$$

Summary

- An **access plan** is a detailed plan for the execution of a query. It describes:
  - the sequence of basic operations (select, project, join, etc.) used to process the query
  - how each operation will be implemented, e.g., which join method will be used, which indices will be used to perform a selection.
- Each operation may be assigned a cost, and an access plan may be assigned a cost by summing the costs of its component operations.
- The cost assigned to an operation depends on assumptions about the data in the relations.